

A Reduction of Finitely Expandable Deep Pushdown Automata

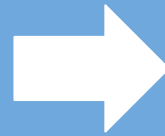
Lucie Charvát and Prof. Alexander Meduna

Introduction

For a positive integer n , n -expandable deep pushdown automata always contain no more than n occurrences of non-input symbols in their pushdowns during any computation. As its main result, the present paper demonstrates that these automata are as powerful as the same automata with only two non-input pushdown symbols $\$$ and $\#$, where $\#$ always appears solely as the pushdown bottom.

Construction

States in Q_R include not only the states corresponding to the states in Q but also strings of non-input symbols. Whenever M pushes a non-input symbol onto the pushdown, M_R records this information within its current state and pushes $\$$ onto the pushdown instead.

$$M = (\{s, q, t\}, \{a, b, c\}, \Gamma, R, s, S, \{f\})$$
$$\Gamma = \{a, b, c, A, S, \#\}$$
$$R = \{1sS \ qAA, 1qA \ fab, 1fA \ fc, 1qA \ taAb, 2tA \ qAc\}$$

$$M_R = (Q_R, \{a, b, c\}, \Gamma_R, R_R, \langle s;S \rangle, \$, F_R)$$
$$\Gamma_R = \{a, b, c, A, S, \#\}$$
$$Q_R = \{\langle s;S \rangle, \langle q;AA \rangle, \langle f;A \rangle, \langle f;E \rangle, \langle t;AA \rangle\}$$
$$F_R = \{\langle f;A \rangle, \langle f;E \rangle\}$$
$$R_R = \{1\langle s;S \rangle \$ \ \langle q;AA \rangle \$ \$, 1\langle q;AA \rangle \$ \ \langle f;A \rangle ab, 1\langle f;A \rangle \$ \ \langle f;E \rangle c, 1\langle q;AA \rangle \$ \ \langle t;AA \rangle a \$ b, 2\langle t;AA \rangle \$ \ \langle q;AA \rangle \$ c\}$$
